# You Don't Need a Ledger: Lightweight Decentralized Consensus Between Mobile Web Clients

SERIAL '19

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## eLoyalty: loyalty points across merchants

Integrated loyalty programs between several small stores Redeem points at any participating store

- Double-spending problem Customer spends same loyalty point twice
- Decentralized: no central authority Merchants do not fully trust each other

## Your local merchant can't run a blockchain node

Expensive hardware requirements

Requires large amount of processing power and storage space

## Complex backend setup

Requiring large infrastructure to setup permissioned network

## Hard to setup consortia between companies

High legal and monetary burden



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## eLoan: loans with unpaid invoices as collateral

Financial institutions offer loans to companies Using unpaid invoices as collateral, verified by the network

- Double-spending problem Company uses invoice twice
- Financial institutions boycott each other They all want to give the loan themselves

## Your bank is not going to rely on proof-of-work

## No consensus finality

Forks can happen with decreasing probability over time

## Blockchain keeps track of all data

Can impose problems with privacy laws such as GDPR

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Architecture and consensus protocol

Evaluation

Zooming out

State-of-the-art and future work

## Lightweight middleware for consensus

Performance and infrastructure

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## Lightweight middleware for consensus Architecture and consensus protocol

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## Use peer-to-peer state-synchronization with voting for Byzantine fault-tolerance

State-based approach using signed data-structures No ledger with history as in operation-based approaches

Peer-to-peer state-synchronization No central component that manages the system

Voting for Byzantine fault-tolerance Tolerate both crash failures as well as malicious nodes

Lightweight setup

Client-side peer-to-peer network running in the browser

# The lightweight architecture requires only two small backend components



# There is no distributed ledger with coins, only tickets which can be used only once



Coins can be used forever

Ownership changes each time it is used



Only track the unspent tickets No need for a ledger

## Protocol has two steps

Issue new token 1.

2.

Redeem token for asset (e.g. loan or loyalty reward)

## Protocol to issue tokens

Customer asks a new token from a node 1. Providing its public key and application-specific info





## Protocol to issue tokens

- Customer asks a new token from a node 1.
- Node verifies the info and creates a new token 2. Stores the token in the local database and sends it back to the customer



## Protocol to issue tokens

- Customer asks a new token from a node 1.
- Node verifies the info and creates a new token 2.
- Node synchronizes token asynchronously with other nodes 3. Using a signature to guarantee authenticity and a timestamp against replay attacks



Customer sends ID of token to node and asks to redeem it 1. Using its private key to prove ownership





- Customer sends ID of token to node and asks to redeem it 1.
- Node verifies token and starts a vote to redeem it 2. A node can only start a vote when customer has sent his signature



- 1. Customer sends ID of token to node and asks to redeem it
- 2. Node verifies token and starts a vote to redeem it
- 3. Other nodes receive votes and vote themselves Every node can vote exactly once for each token



- 1. Customer sends ID of token to node and asks to redeem it
- 2. Node verifies token and starts a vote to redeem it
- 3. Other nodes receive votes and vote themselves
- 4. Original node waits until it receives  $\frac{2}{3} \times n + 1$  of the votes Gives the real-life product to the customer afterwards



# The protocol guarantees liveness and safety for honest customers

Safety: nothing bad happens

No double spending of tokens

Liveness: eventually something good happens Tokens can actually be redeemed

Malicious customers can lock their tokens and lose them But safety is still guaranteed

## Adversary model

- Tolerates up to  $\frac{1}{3} \times (n 1)$  byzantine nodes Tolerating both crash failures and malicious nodes



- Attacker cannot break the used cryptographic primitives ECDSA with P-256 and SHA-256
- **P2P** discovery service must be honest Otherwise liveness will be broken due to network partitions

# Designed for a closed group of members with infrequent changes

Members can be added or removed using the same protocol But with manual voting and verification

New members only participate in votes for new tokens Membership is fixed on token creation

If too many members leave, older tokens become invalid The <sup>2</sup>/<sub>3</sub> majority for those tokens cannot be reached anymore

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## The middleware scales up to 50 nodes



## Each node can have 1000s of end-users

eLoyalty: nodes are the merchants

Scale: local farmer's market or shopping street

## eLoan: nodes are the financial institutions

Scale: small country with 50 different banks

## The lightweight middleware requires only two small backend components

This middleware

Web server, signaling server

Hyperledger Fabric using Hyperledger composer

Web server, peers, orderers, REST-, CouchDB- and CA-servers, membership service

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## State-of-the-art and related work

Peer-to-peer data synchronization systems Legion, Yjs

PoW blockchains

Bitcoin, Ethereum

Blockchains using BFT consensus

Hyperledger Fabric, HoneyBadger, Algorand, Avalanche

Off-chain protocols

Bitcoin lightning network

## The middleware solves the double-spending problem in a specific range of applications

Keep track of transient resources Not for perpetual resources such as coins

Un-frequent membership changes Members can leave and join at a slow rate

Up to 50 nodes that are online most of the time Network can be unstable, some nodes may be offline

## Challenges and future work

One ECDSA signature per token per node for consensus

Not scalable to 100-1000 nodes

Look into non-interactive signature aggregation schemes E.g. BLS

Current P2P network is constructed randomly

Look into structured networks such as fat-tree overlay



## You Don't Need a Ledger:

Only keep the current state using signed data-structures

# Lightweight Decentralized Consensus

Simple voting for asynchronous BFT between the clients

# Between Mobile Web Clients

Running in a web browser on a mobile device

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